

# APPENDIX

## A. Algorithms

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**Algorithm 1** Conversion of Wireless Data Broadcast into DAG (CWDB-DAG)

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**Input:** Data items to be broadcast on the  $i$ th channel in the  $j$ th time slot  $c(i, j)$ ,  $1 \leq i \leq N, T_1 \leq j \leq T_2$ ;  
**Output:** Converted DAG  $G(V, E)$ , it is denoted by Matrix  $M(num, num)$ ;

```
1: Let  $num = 0$ ;  
2: for  $i = 1$  to  $N$  do  
3:   for  $j = T_1$  to  $T_2$  do  
4:     if  $c(i, j) \neq NULL$  then  
5:        $num++$ ;  $node_{num}.time = j$ ;  $node_{num}.data = c(i, j)$ ;  $node_{num}.channel = i$ ;  
6:     end if  
7:   end for  
8: end for  
9:  $M(num, num) = 0$ ;  
10: for  $i = 1$  to  $num$  do  
11:   for  $j = i + 1$  to  $num$  do  
12:     if  $node_i.data \neq node_j.data$  and  $node_i.channel = node_j.channel$  then  
13:       if  $node_i.time < node_j.time$  then  
14:          $M(i, j) = 1$ ;  $node_j.indegree++$ ;  $node_i.outdegree++$ ;  
15:       break;  
16:     else if  $node_i.time < node_j.time$  then  
17:        $M(j, i) = 1$ ;  $node_i.indegree++$ ;  $node_j.outdegree++$ ; break;  
18:     end if  
19:   end if  
20: end for  
21: end for  
22: for  $i = 1$  to  $num$  do  
23:   for  $j = 1$  to  $num$  do  
24:     if  $node_i.data \neq node_j.data$  and  $node_i.channel \neq node_j.channel$  then  
25:       if  $node_i.time < node_j.time - 1$  then  
26:          $M(i, j) = 1$ ;  $node_j.indegree++$ ;  $node_i.outdegree++$ ; break;  
27:       else if  $node_j.time < node_i.time - 1$  then  
28:          $M(j, i) = 1$ ;  $node_i.indegree++$ ;  $node_j.outdegree++$ ; break;  
29:       end if  
30:     end if  
31:   end for  
32: end for
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**Algorithm 2** Discovery of All Paths (DAP)

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**Input:**  $M(num, num)$ ;  
**Output:** set of all paths  $P$ ;  
1:  $P \leftarrow NULL$ ;  
2: find the node set  $V_{in}$  which indegree of node is 0;  
3: find the node set  $V_{out}$  which outdegree of node is 0;  
4: **while**  $V_{in} \neq NULL$  **do**  
5:   get a node  $v_{in_i} \in V_{in}$ ;  
6:   get a node  $v_{out_j} \in V_{out}$ ;  
7:   find a path  $p_k$  from  $v_{in_i}$  to  $v_{out_j}$  through using DFS;  
8:   mark all edges in  $p_k$ ;  
9:   **if** all edges  $e \in E$  are marked **then**  
10:    **break**;  
11:   **else**  
12:     $P \leftarrow p_k$ ;  
13:   **end if**  
14: **end while**

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**Algorithm 3** Approximate Data Retrieval on Single Antenna (ADR-SA)

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**Input:** broadcast cycle  $[T_1, T_2]$ , channel set with requested data items  $C$ , the set of requested data items  $D_k$ , and the number of channels  $N$ ;  
**Output:** optimal data retrieval sequence  $P_{opt}$ ;  
1:  $t = T_1$ ;  
2: **while**  $t \leq T_2$  **do**  
3:    $converted\ DAG(C, T_1, T_2, N)$ ;  
4: **end while**  
5:  $DAP(M)$ ;  
6: find the minimal set cover  $S$  based on  $P$  and  $D_k$ ;  
7: according to  $D_k$ , delete un-requested data items in  $S$ ;  
8:  $P_{opt} \leftarrow S$ ;

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**Algorithm 4** Approximate Data Retrieval on Multiple Antennae (ADR-MA)

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**Input:** broadcast cycle  $[T_1, T_2]$ , channel set with data items  $C$ , the set of requested data items  $D_k$ , and the number of channels  $N$ ;

**Output:** optimal data retrieval sequence  $P_{opt_1}, P_{opt_2}$ ;

```
1:  $t = T_1, antennae = 2$ ;  
2: while  $t \leq T_2$  do  
3:    $converted\ DAG(C, T_1, T_2, N)$ ;  
4: end while  
5:  $Path = DAP(M)$ ;  
6: if  $|D_k| \leq a$  then  
7:    $D_k$  is divided into two subsets of data items  $D_{k_1}, D_{k_2}$ , and store all patterns in  
    $PAT = \{pat_1, pat_2, \dots, pat_{|PAT|}\}$ ;  
8: else  
9:   for  $i = 1$  to  $\lfloor \frac{|D_k|}{2} \rfloor$  do  
10:    randomly select  $i$  patterns  $pat$  of the split  $\{D_{k_1}, D_{k_2}\}$  on  $D_k$  where  $|D_{k_1}| = i$   
    and  $|D_{k_2}| = |D_k| - i$ , and store them to  $PAT_i$ ;  
11:     $PAT \leftarrow PAT \cup PAT_i$ ;  
12:   end for  
13: end if  
14: for  $i = 1$  to  $|PAT|$  do  
15:   find the minimal set cover sequences  $P_{opt_1}$  and  $P_{opt_2}$  from  $Path$  on the basis of  $D_{k_1}$  and  
    $D_{k_2}$  in  $pat_i$ ;  
16:   according to  $P_{opt_1}$  and  $P_{opt_2}$ , compute the access latency and store as  $t_{pat_i}$ ;  
17: end for  
18: find a pattern which has the minimal access latency and output the corresponding retrieval  
   sequences on two antennae  $P_{opt_1}$  and  $P_{opt_2}$ .
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## B. Experimental Results

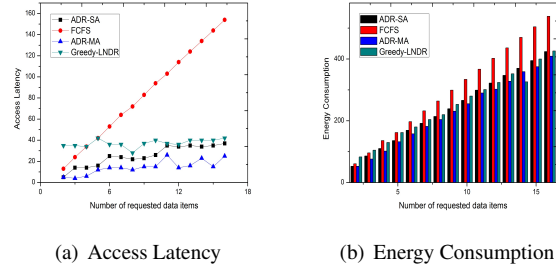


FIGURE 7  
Experimental results of access latency and energy consumption in wireless data broadcast with four channels

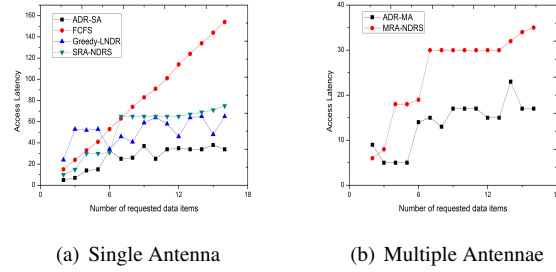


FIGURE 8  
Experimental results of access latency in wireless data broadcast with six channels and the client on single antennae and multiple antennae

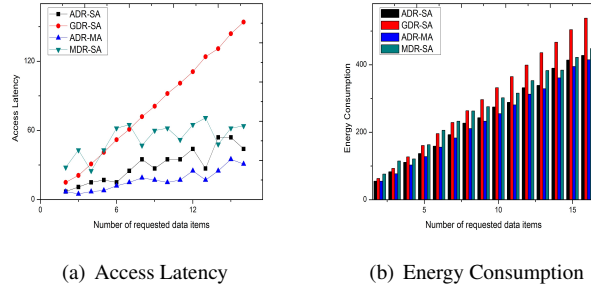


FIGURE 9  
Experimental results of access latency and energy consumption in wireless data broadcast with eight channels

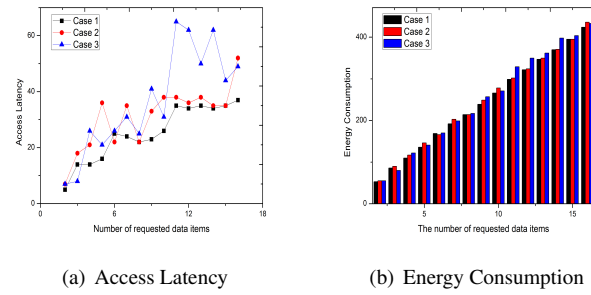
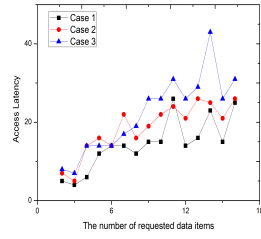
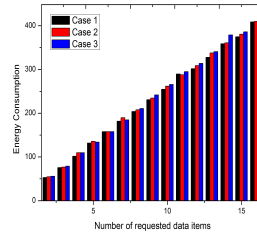


FIGURE 10  
Experimental results of access latency and energy consumption for ADR-SA applied to three cases in wireless data broadcast with four channels



(a) Access Latency



(b) Energy Consumption

FIGURE 11  
Experimental results of access latency and energy consumption for ADR-MA applied to three cases in wireless data broadcast with four channels